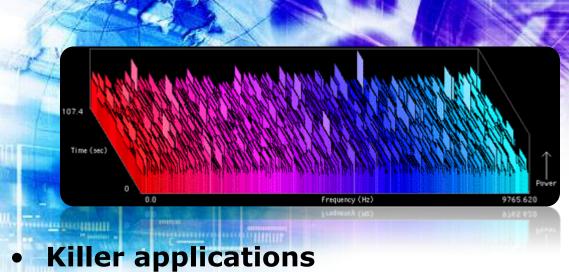
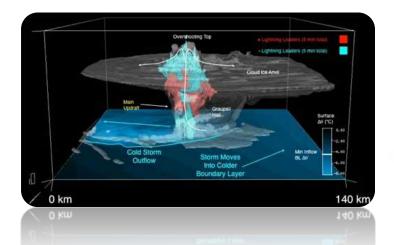
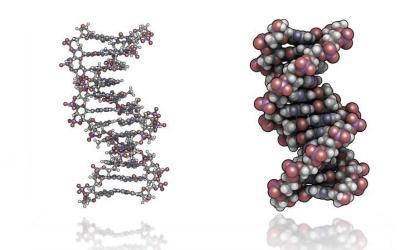
Parallel Applications Design with MPI

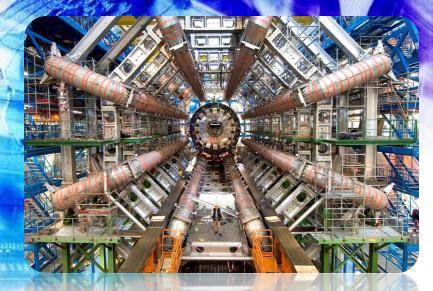


Science Research Challanges

- Challenging use of computer power and storage
- Who might be interested in those applications?
- Simulation and analysis in modern science







Example: Large Hadron Collider as (CERN)

LHC Computing Grid

- Worldwide collaboration of > 170 computing centers in 34 countries
- Recording rate (raw) 1 GByte/sec
- Sum between 5 up to 8 PetaByte/year (10¹⁵ B/year)
- Estimated to be 200,000 faster than some today's fastest CPUs
- How can we get this working?



The need for HPC

Moore's law

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- In 1965 stated that transistor density double at each 1.5 years

- A wrong assumption is that if transistor density double than computer gets twice as fast every 18/24 months

 This is wrong: the proof is that today highest clocks don't reach 20GHz



Can we use a single processor?

- Short answer is: no
 - Curiously Moore's law is still true
 - Transistor densisty indeed doubles but higher clock rate leads to unmanageable heat and power consumption
 - Super computers are too expensive for medium size problems
 - The solution is work with multiple processors at the same time
 - If super computers are too expensive why not create a machine clustering desktop solutions

Apr-98 Feb-99 Dec-99 Oct-00 Aug-01 Jun-02 Feb-04 Dec-04 Jun-07 Apr-08 Feb-09 Feb-09 Feb-09

Computing

Clusters

Cluster of workstations

- Commodity PCs interconnected through a network
- More affordable than supercomputers
- Adaptable to any scale of usage
- Fast aquired popularity among researchers

• Timeline

 $\mathbf{0}$

- 1993: Beowulf project
- 1997: Berkley NoW is first cluster on top500
- 2010: 80 % of machines in top500 are clusters
- Source: www.top500.org

Programming on a Cluster

Applications must be rewritten

- Proccess communication changes
 - Instead of memory must use the network

Possible solutions

-Ad Hoc

- Work only for the platform it was designed for

-PVM

Research project for heterogeneous network computing

- MPI
 - It's a standard, independent of implementation
 - Have more than three free implementations
 - Here we are going to talk about MPI

MPI – Message Passing Interface

MPI in a nutshell

- It is a library specification

- Works natively with C and Fortran
- Not a specific implementation or product
- -Scalable
 - Must handle multiple machines
- -Portable
 - Sockets API change from one OS to another
 - Handles Big-endian/little-endian architectures
- Efficient
 - Optimized communication algorithms
 - Allow communication and computation overlap

MPI – Message Passing Interface

MPI References

Books

– Using MPI: Portable Parallel Programming with the Message Passing Interface, by Gropp, Lusk, and Skejellum, MIT Press, 1994.

– MPI: The Complete Reference, by Snir, Otto, Huss-Lederman, Walker, and Dongarra, MIT Press, 1996.

 Parallel Programming with MPI, by Peter Pacheco, Morgan Kaufmann, 1997.

- The standard:

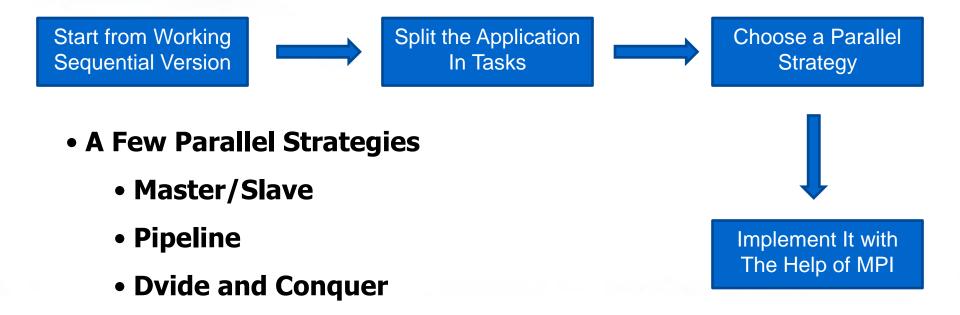
- at http://www.mpi-forum.org

• MPI

- Use of a single program, on multiple data
- What does it do?
 - way of identifying process
 - Low-level independent API
 - Optimized communication
 - Allow communication and computation overlap
- What does it do not?
 - gain performance of application for free
 - application must be adapted

Possible Programming Workflow

COLUMN 111



Master/Slave

TRACE OF BRIDE

• Master is one process that centrilizes all tasks

•Slaves starve for work

.....



- Master/Slave
 - Master is often the bottleneck
 - Scalability is limited due to centralization
 - Possible to use replication to improve performance
 - It is adatable to heterogenous platforms

Task 1

Task 2

Task 3

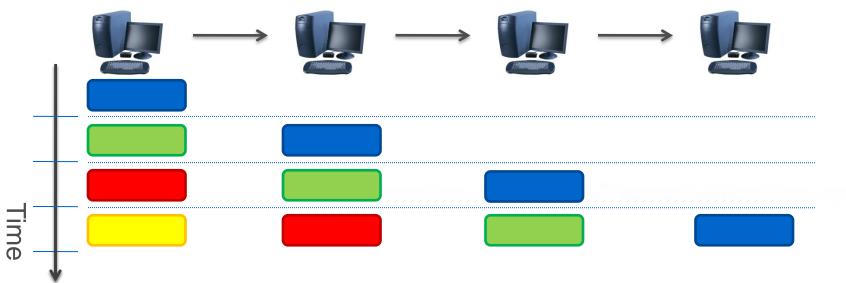
Task 4

- Each process plays a specific role, pipeline stages
- Data follows in a single direction

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Pipeline

• Parallelism is achieved when the pipeline is full



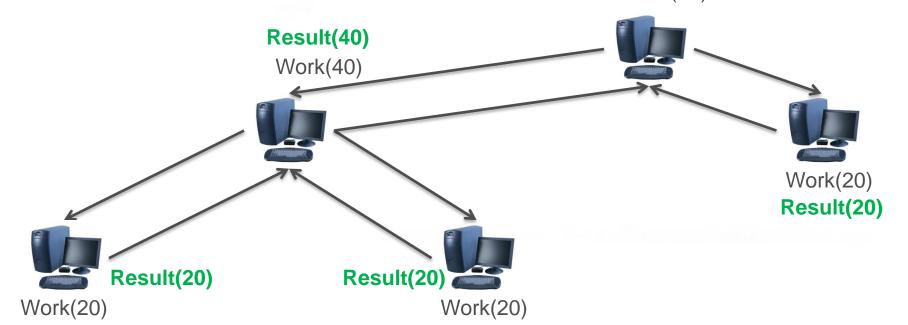


- Scalabillity is limited by the number of stages
- Synchronization may lead to bubbles
 - Slow sender
 - Fast receiver
- Difficult to use on heterogenous platforms

Divide and Conquer

- Recursevely partion task on roughly equal sized tasks
- Or process the taks if it is small

Result(60) Work(60)



- Divide and Conquer
 - More scalable
 - Possible to use replicated branches
 - In practice is difficult to split tasks
 - Suitable for branch and bound algorithms

• Installing

CALCULATER FOR

- Some common MPI implementations, all free:
 - OpenMPI

http://www.open-mpi.org/

– MPICH-2
http://www.mcs.anl.gov/research/projects/mpich2/

- LAM/MPI

http://www.lam-mpi.org/

- Installing
 - I'm using MPICH-2
 - Installed in Ubuntu 10.04 Lucid Lynx with
 - \$ sudo apt-get install mpich2
 - Should work for most Debian based distributions
 - Must create a local configuration file
 - \$ echo "MPD_SECRET_WORD=ChangeMe" > ~/.mpd.conf

Test program

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#include <mpi.h>
#include <stdio.h>

int main(int argc, char **argv){

```
/* Initialize MPI */
MPI_Init(&argc, &argv);
```

printf("Test Program\n");

```
/* Finalize MPI */
return MPI_Finalize();
}
```

Compiling

COLUMN AND ADDRESS

 Compiled with gcc, but a mpicc script is provided to invoke gcc with specific MPI options enabled

\$ mpicc mpi_program.c -o my_mpi_executable

- Executed with a specital script

- \$ mpirun -np 1 my_mpi_executable
- \$ mpirun -np 2 my_mpi_executable
- \$ mpirun -np 3 my_mpi_executable

Running

 Compiled with gcc, but a mpicc wrapper is provided to invoke gcc with specific mpi options

- \$ mpicc mpi_program.c -o my_mpi_executable
- For a complete list of parameters try
 - \$ man mpicc
- Executed with a special wrapper
 - \$ mpirun -np 2 my_mpi_executable

• Exercise 1 – Hello World

• Compile and run the simplest MPI program that only prints the "Hello World" string and after exits

• Try vary the -np <nproc> parameter and observe the differences

How many proccess are running?

int MPI_Comm_size(MPI_Comm comm, int *psize)

- comm

- Group of process to communicate
- For grouping all process use MPI_COMM_WORLD

- psize

– Passed as reference will return the total amoung of proccess in this communicator

• Exercise 2 – Number of Proccess

• Create program that prints the total number of available process on the screen

 Vary the "-np <param>" to verify that your program is working

Assigning Process Roles

int MPI_Comm_rank(MPI_Comm comm, int *rank)

- comm

- Group of process to communicate
- To group all available process use MPI_COMM_WORLD

- rank

– Passed as reference will return the unique ID of the calling process in this communicator

- Exercise 3 Who am I?
 - •If I am process 0
 - Prints: "hello world"
 - else

COLUMN AND ADDRESS

- Prints: "I'm process <ID>"
- Replacing <ID> by the process rank

Running on Grid5000 1/2

- Log in grid5000 using the instructions gave to you
 - \$ ssh <username>@access.<frontend>.grid5000.fr

-Log in the front end of your choice

- \$ ssh <frontend>
- Make a reservation
 - \$ oarsub -1 nodes=4,walltime=1 -I

Running on Grid5000 2/2

 When connected to a cluster in interactive mode a file with the list o available machines is generated

- \$ cat \$OAR_FILENODES
- Compile your code again
 - \$ mpicc -o mybin myprogram.c
- Run MPI application

\$ mpirun --mca plm_rsh_agent oarsh -machinefile
\$OAR_FILENODES -np <nproc> mybin

• Exercise 4 – Running on a cluster

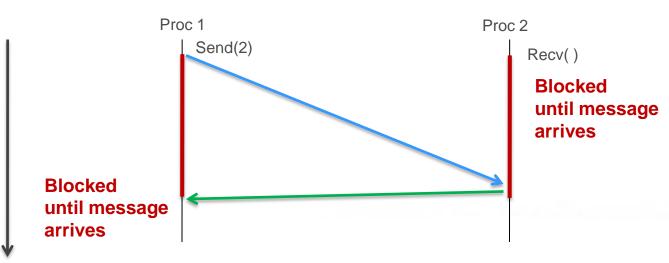
• Use your Grid5000 account to run the "Who am I?" program in a real cluster

• Use the gethostname function to print the host name where the process is running

MPI One-to-one Communication

Synchronous/Blocking

- Process sits waiting for message to arrive
- Synchronization purpose



Time

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Blocking Send 1/2

int MPI_Send(void *buf, int count, MPI_Datatype dtype,

int dest, int tag, MPI_Comm comm)

- buf

- Pointer to the data to be sent

- count

- Number of data elements in buf
- dtype
 - Type of elements in buf

Blocking Send 2/2

int MPI_Send(void *buf, int count, MPI_Datatype dtype,

int dest, int tag, MPI_Comm comm)

- dest
 - Rank of destination process
- tag
 - Tag another integer to identify the message
- comm
 - Same as before, for all proccess use MPI_COMM_WORLD

Blocking Receive 1/2

int MPI_Recv(void *buf, int count, MPI_Datatype dtype,

int src, int tag, MPI_Comm comm, MPI_Status &status)

- buf

- Pointer where data will be received if succeed

- count

- Maximum number of elements that buf can handle

- dtype
 - Type of elements in buf

Blocking Receive 2/2

int MPI_Recv(void *buf, int count, MPI_Datatype dtype,

int src, int tag, MPI_Comm comm, MPI_Status &status)

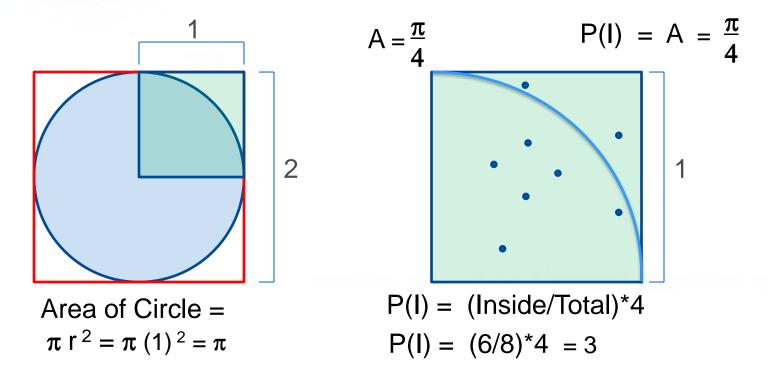
- src

- Rank of sender process
- tag
 - Message tag
- stat

 Sending process info, if desired can be ignored using MPI_STATUS_IGNORE

• Exercise 5 (1/2) – Computing π by Monte Carlo Methods

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- Exercise 5 (2/2) Computing π by Monte Carlo Methods
 - Generate two random numbers X, Y in [0,1]
 - If (X*X + Y*Y) <= 1
 - Add 1 to counter
 - At the end use (counter/total)*4 as π approximation
 - \bullet More random points generated more close to π

- Performance Evaluation
 - Elapsed Time
 - The timer itself

- Speedup

COLUMN THE R

How many times my application is faster than the sequential version

- Efficiency

– Estimate processing power dissipation

#include <mpi.h>
#include <stdio.h>

int main(int argc, char **argv){

```
/* Initialize MPI */
MPI_Init(&argc, &argv);
```

```
t1 = MPI_WTime();
compute_pi();
t2 = MPI_WTime();
printf("Elapsed time: %f\n", t2 - t1);
```

```
/* Finalize MPI */
return MPI_Finalize();
}
```

double MPI_WTime()

- RETURN

Timer

The time passed,
 in seconds, since
 an arbitrary time
 in the past

- Performance Evaluation
 - Speedup
 - Obtained from elapsed time
 - Ratio of elapsed time with one processor and elapsed time with n processors

$- Speedup(n) = \frac{T(1)}{T(n)}$

- -T(1) = elapsed time with one processor
- -T(n) = elapsed time with n processors
- The ideal is: Speedup(i) = i
- Meaning: using i processors I get i times faster

- Performance Evaluation
 - Efficiency
 - It is obtained from speedup
 - The efficiency shows the percentage of usage by processor

$- Efficiency(n) = \frac{Speedup(n)}{n}$

- Ideal is: Efficiency(i) = 1
- Meaning:

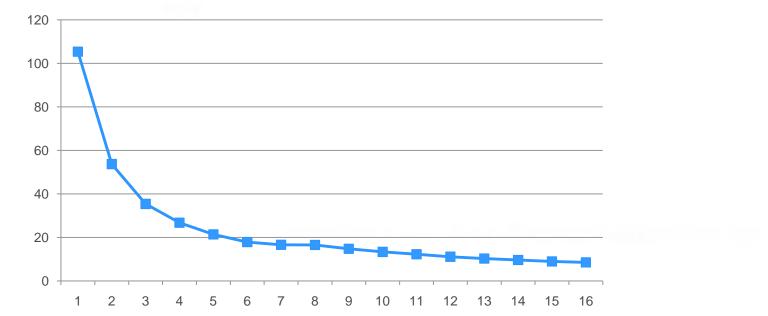
Each processor is being used at 100%, no parallelization overhead

- Exercise 6 Performance Evluation of the π Problem
 - Compute the elapsed time to compute 10⁹ iterations of Pi with:
 - 1, 2, 4, 8, 10 processors
 - Present speedup and efficiency

• Performance Evluation of the π Problem

• Elapsed Time

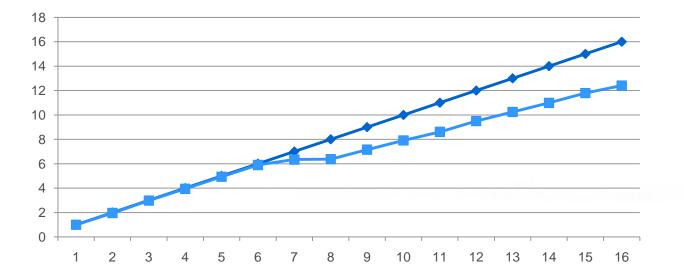
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- Performance Evluation of the π Problem
 - Speedup

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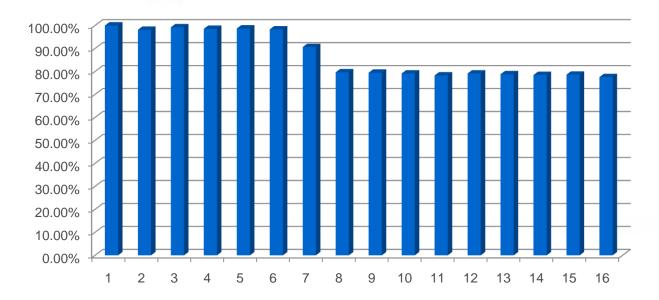




• Performance Evluation of the π Problem

• Efficiency

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Deadlock

COLUMN 11

- Two process, or more, are blocked waiting for the other

-The computation do not advance

Deadlock

```
if(my_Rank == 0){
    MPI_Recv(&tmp, 1, MPI_INT, 1, 122, MPI_COMM_WORLD, MPI_IGNORE_STATUS);
    MPI_Send(&dat, 1, MPI_INT, 1, 122, MPI_COMM_WORLD);
}
if(my_Rank == 1){
    MPI_Recv(&tmp, 1, MPI_INT, 0, 122, MPI_COMM_WORLD, MPI_IGNORE_STATUS);
    MPI_Send(&dat, 1, MPI_INT, 0, 122, MPI_COMM_WORLD);
}
```

Deadlock

• Can be solved proper ordering blocking calls

```
if(my_Rank == 0){
    MPI_Recv(&tmp, 1, MPI_INT, 1, 122, MPI_COMM_WORLD, MPI_IGNORE_STATUS);
    MPI_Send(&dat, 1, MPI_INT, 1, 122, MPI_COMM_WORLD);
}
if(my_Rank == 1){
    MPI_Send(&dat, 1, MPI_INT, 0, 122, MPI_COMM_WORLD);
    MPI_Recv(&tmp, 1, MPI_INT, 0, 122, MPI_COMM_WORLD, MPI_IGNORE_STATUS);
}
```

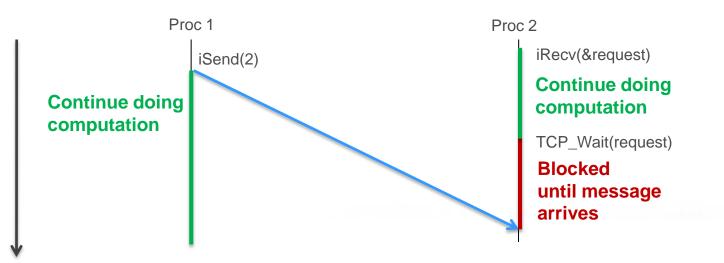
MPI One-to-one Communication

Assynchronous/Non-Blocking Receive

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Time

- Process tries to receive, returns if there is no message
- Wait for the message is done using Wait



Non-blocking Send and Receive

```
int MPI_Isend(..., MPI_Request &req)
```

- req

Reference to the communication request, holds information to use later

int MPI_Irecv(..., MPI_Request &req, MPI_Status &status)

- req

Reference to the communication request, holds information to use later

• MPI_Wait

int MPI_Wait(MPI_Status *status, MPI_Request *req)

- status

 Contains information about the received message can be ignored using MPI_STATUS_IGNORE

– req

– Reference to the communication request, holds information to use later

MPI_Waitany

int MPI_Waitany(int count, MPI_Status *status[], int *index, MPI_Request *req[])

- count

– Number process which are going to wait for

- status

– Array of statusus, to ignore use MPI_STATUS_IGNORE

- index

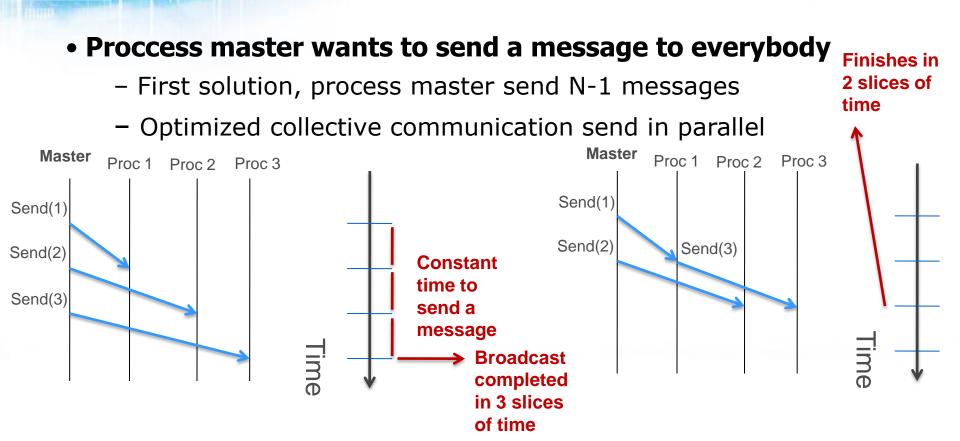
- Returns the index of the received request

- req

- Array of requests to wait for

- Exercise 7 – π Problem with MPI_Waitany
 - Use MPI_Isend and MPI_Irecv instead of blocking communication
 - Use MPI_Waitany to receive messages efficiently

MPI Collective Communication



TRACE OF BRIDE

• Broadcast

int MPI_Bcast(void* buffer, int count, MPI_Datatype
datatype, int root , MPI_Comm coom)

- root

– If (myrank == root) send the content of buffer else not receive the content through buffer parameter

– It is the first process to send messages to the others

• Exercise 8 – π Problem with MPI_Bcast

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Collective reduce 1/2

int MPI_Reduce(void* *sendbuf, void *recvbuf, int count, MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm)

- sendbuf

- Data to be sent
- recvbuf
 - Data to be received
- count
 - Number of elements in sendbuf and recvbuf
- datatype
 - Datatype of sendbuf and recvbuf elements

Collective reduce 2/2

int MPI_Reduce(void* *sendbuf, void *recvbuf, int count, MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm)

- MPI_Op

– The arithmetic operation to execute some possible values: MPI_MAX, MPI_MIN, MPI_SUM, MPI_PROD, and so on STATUS_IGNORE

- root

– The same meaning as for MPI_Bcast

- Exercise 9 – π Problem with MPI_Reduce
 - Use MPI_Isend and MPI_Irecv instead of blocking communication
 - Use MPI_Waitany to receive messages in the efficientest order

Conclusion

- MPI is useful and easy to program
- Well, at least easyer than sockets
- Have many vendor and free implementations
- It is optimized for different network architectures
- What we saw:

– Some of the MPI one-to-one and collective communication API

- What we didn't saw:

– A lot of stuff, MPI_Scatter, MPI_Barrier, MPI_Get_processor_name, and so on...

• Future of MPI

ACCOUNT OF A DESIGNATION OF

- Exascale computing??
- 10^18 Flops/s

